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BTREE: A FORTRAN CODE FOR B+ TREE(U) NAVAL SURFACE
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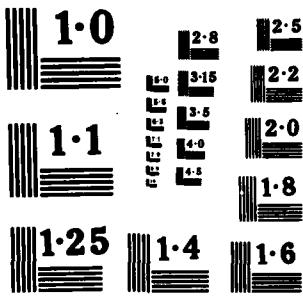
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REPORT DOCUMENTATION PAGE		READ INSTRUCTIONS BEFORE COMPLETING FORM
1. REPORT NUMBER NSWC TR 85-54	2. GOVT ACCESSION NO. AD-A157	3. RECIPIENT'S CATALOG NUMBER 026
4. TITLE (and Subtitle) BTREE: A FORTRAN CODE FOR A B+ TREE		5. TYPE OF REPORT & PERIOD COVERED Final: Fiscal Year '85
7. AUTHOR(s) Elliot Winston		6. PERFORMING ORG. REPORT NUMBER
9. PERFORMING ORGANIZATION NAME AND ADDRESS Naval Surface Weapons Center (R44) 10901 New Hampshire Avenue Silver Spring, MD 20903-5000		10. PROGRAM ELEMENT, PROJECT, TASK AREA & WORK UNIT NUMBERS 64601N; S0267; 0; SU150D
11. CONTROLLING OFFICE NAME AND ADDRESS		12. REPORT DATE 1 April 1985
		13. NUMBER OF PAGES 63
14. MONITORING AGENCY NAME & ADDRESS (if different from Controlling Office)		15. SECURITY CLASS. (of this report) UNCLASSIFIED
		15a. DECLASSIFICATION/DOWNGRADING SCHEDULE
16. DISTRIBUTION STATEMENT (of this Report) Approved for public release, distribution unlimited		
17. DISTRIBUTION STATEMENT (of the abstract entered in Block 20, if different from Report)		
18. SUPPLEMENTARY NOTES		
19. KEY WORDS (Continue on reverse side if necessary and identify by block number) B+ Tree, Database Manager, Node, Leaf, Root		
20. ABSTRACT (Continue on reverse side if necessary and identify by block number) This report discusses and documents a FORTRAN code for a B+ Tree, a data structure which if frequently used is the foundation of a database manager.		

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FOREWORD

This report contains documentation for a FORTRAN implementation of a B+ tree, a data structure which is often used as the foundation of a database manager. Because the code is written in a high-level language, it is basically transportable to any computer with FORTRAN capability (minor modification may be required for compatibility with a host computer's operating system and compiler). The work was done as a first step towards developing a user-friendly, interactive database manager needed by US1 to support studies requiring the extensive use of minefield planning codes.

This work has been supported by the Mine Improvement Program at NSWC under Project S0267.

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Additional keywords: Subroutine documentation
Specification: —



CONTENTS

	<u>Page</u>
INTRODUCTION.....	1
SPECIFICATIONS.....	2
USAGE.....	2
RECONFIGURATION.....	3
REFERENCES.....	5
APPENDIX A - SUBROUTINE DOCUMENTATION.....	A-1
APPENDIX B - FORTRAN CODE LISTING.....	B-1
DISTRIBUTION.....	(1)

ILLUSTRATION

<u>Figure</u>	<u>Page</u>
1 EXAMPLE OF A B+ TREE.....	1

INTRODUCTION

A B+ tree is a data structure which is particularly well suited for storing the keys which identify the records in a database. The objective of a B+ tree, hereafter referred to simply as a tree, is to minimize the number of mass storage accesses required to find a specified key. A conceptual representation of a tree, containing the letters {B,D,E,G,H,L,N,R,S} as keys, is shown in Figure 1.

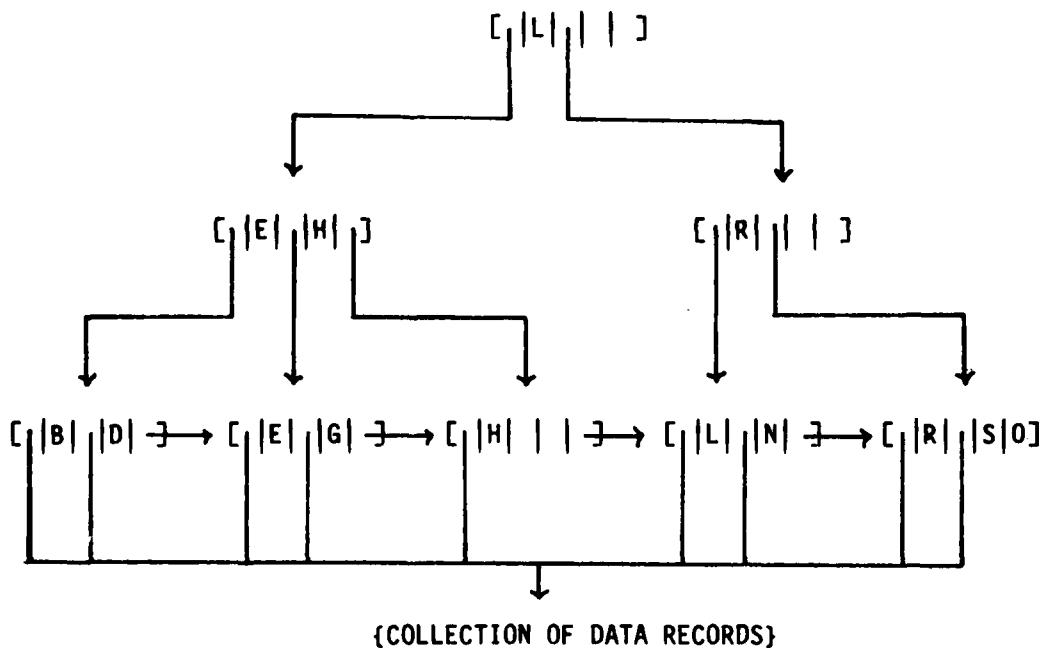


FIGURE 1. EXAMPLE OF A B+ TREE

A tree corresponds to a file, each node of the tree corresponds to a record in that file, and a pointer between nodes corresponds to the record number of the node to which it points. The tree in Figure 1 has a nodal capacity equal to 2. The nodes on the bottom level, called the leaves of the tree, contain all the keys. Each leaf points to its neighbor on the right, and the rightmost leaf points to zero, indicating that it is the last leaf in the tree.

All key searches begin at the top node, known as the root. To find the letter G, for example, the left node on the second level is searched after the root because G precedes L in alphabetical order. Since G is between E and H, the second leaf from the left is searched next, and the key is found. Thus, the number of accesses required to find a key is equal to one more than the height of the tree. As keys are added to the tree, leaves become full and split in half; as keys are removed, adjacent leaves may merge. A complete description of splitting and merging rules will not be given here, but the interested reader is referred to the excellent introduction presented by Comer¹; a more analytical discussion can be found in Knuth².

Documentation for a database manager based on the code listed in Appendix B is contained in Winston³.

SPECIFICATIONS

The implementation listed in Appendix B is written in a version of FORTRAN 77 for a DEC VAX/780 computer with the VMS operating system. In particular, file names are at most 9 characters long, and have extenders with as many as 3 characters. A tree called (name) can have as many as 3 files associated with it: (name).KEY, (name).NOD, and (name).REC. The file (name).KEY corresponds to the tree itself, and (name).NOD and (name).REC are node and record stacks, respectively. The numbers of the nodes and data records that are deleted from the tree as a result of key deletions are saved by the stacks and reused as needed. Stack files containing no numbers are automatically erased from the system.

Records corresponding to nodes have a length of 256 bytes, and the maximum tree height is equal to 5. Keys have a maximum length of 20 characters, and the tree can accommodate up to 65,535 keys. Application programs which call BTREE can have up to 10 trees open simultaneously, via the logical device unit numbers 1,2,...,10.

USAGE

A call to BTREE is accomplished by the standard FORTRAN syntax

```
CALL BTREE (LTR,LDU,A,MAXLEN,IREC,IERR).
```

A description of the input and output parameters follows:

INPUT:

LTR (CHARACTER*1)	COMMAND
A -----	(A)dd a key to the tree
C -----	(C)reate a new tree
D -----	(D)elete a key from the tree
F -----	(F)irst key in the tree
G -----	(G)et first occurrence of a partial key
O -----	(O)pen an old tree
S -----	(S)uccessor key
LDU (BYTE)	unit number under which the tree communicates with mass storage; permissible values are 1,2,...,10
A (CHARACTER*20)	full key (required for LTR = 'A', 'D') partial key (required for LTR = 'G') tree name (required for LTR = 'C', 'O')

MAXLEN (INTEGER*4) maximum key length; cannot exceed 20
 (required for LTR = 'C')

OUTPUT:

A (CHARACTER*20)	value of the last key accessed
IREC (INTEGER*4)	number of the data record associated with the last key accessed
IERR (BYTE)	ERROR CODE
0 -----	successful execution of the command
1 -----	illegal value of the parameter 'LTR'
2 -----	attempt to create an existent tree
3 -----	attempt to access a nonexistent tree
4 -----	attempt to find a nonexistent key
5 -----	no successor key exists
6 -----	attempt to insert a key currently in tree

The "get" command is of particular interest because it searches the tree for the first occurrence of a left-justified string within a key. This feature is useful when it is desired to access a particular class of keys, all of which begin with the same string of characters. Setting A equal to the string and executing a "get" command finds the first key which has A as its prefix; successive "successor" commands find the remaining keys in the class. The full key value returned by A should always be checked for its prefix after each such call. This procedure is especially fast because the "successor" command usually does not require a mass storage access. Unlike the "get" command, the "delete" command fails to execute unless the full key value is specified.

Either an "open" command or a "create" command must be executed prior to performing any other operations on a particular tree.

BTREE creates trees which contain unique keys only; an attempt to insert a duplicate key will not execute (IERR = 6). Also, the code prevents the creation of a new tree having the same name as a tree currently on the system (IERR = 2).

RECONFIGURATION

It may be desirable to alter some of the specification parameters to make BTREE mesh more efficiently with a particular application program. These adjustments are indicated as follows:

- (a) To change the maximum key length to k characters, declare the passing parameter A and the internal variable KEYVAL as CHARACTER*k;
- (b) To change the height of the tree to h, declare the arrays BUF(0:h) and PATH(0:h);

- (c) To change the number of trees open simultaneously to t, declare the arrays HTREC(t), MAXREC(t), HTNOD(t), MAXNOD(t), ROOT(t), MKL(t), HEIGHT(t), NAME(t), MARKS(t), MANYS(t), and BLOCS(t);
- (d) To change the size of the nodes to s bytes, change all CHARACTER*256 declarations to CHARACTER*s, assign FULL = (s-4)/ONE in SUBROUTINE BTREE, and assign RECL = s in the OPEN statement in SUBROUTINE NEWTREE.

While it is possible to increase the maximum number of keys allowed in the tree, the alterations required by the code are far more intricate and complex. Changing to a 3-BYTE symbol code for integers would allow up to 16,777,215 keys, but would entail, among other things, rewriting the functions VAL and SYM and checking all sections of the code dealing with loading information into the nodes because 3 bytes must be reserved for each integer instead of 2 bytes. In short, increasing the number of keys is not recommended. Moreover, specially tailored database programs are usually developed, or purchased, to maintain such large databases.

REFERENCES

1. Comer, D., "The Ubiquitous B Tree," Computing Surveys, Vol. 11, 1979, pp. 121-137.
2. Knuth, D., The Art of Computer Programming, Vol. 3: Sorting and Searching (Reading: Addison-Wesley, 1973), pp. 473-480.
3. Winston, E., BOSS: A FORTRAN Code for a Relational Database Manager, NSWC TR 85-56, June 1985.

NSWC TR 85-54

APPENDIX A
SUBROUTINE DOCUMENTATION

SUBROUTINE BTREE

PURPOSE: To control the logic needed to execute the command requested by the value of LTR.

INPUTS:

LTR (= ABC)	CHARACTER*1	command letter
LDU (= IO)	BYTE	input/output device number
A (= NAME) A (= KEYVAL)	CHARACTER*20 CHARACTER*20	key name if LTR = 'A', 'D', or 'G' tree name if LTR = 'C' or 'O'
MAXLEN (= MKL)	INTEGER*4	maximum key length

OUTPUTS:

IREC (= PTR)	INTEGER*4	number of the data record associated with the last key examined
IERR (= ERR)	BYTE	error code number
A	CHARACTER*20	full value of last key examined
ONE	INTEGER*4	parameter equal to MAXLEN + 2
FULL	INTEGER*4	maximum number of keys in a node

EXTERNALS:

ADDKEY, NEWTREE, DELKEY, FIRST, GETKEY, OLDTREE, SUCCESSOR

SUBROUTINE MERGE

PURPOSE: To merge two adjacent nodes into a single node.

INPUTS:

LEAF	LOGICAL	.TRUE. if and only if the current node is a leaf
LFM	INTEGER*4	number of keys in the left node
RTM	INTEGER*4	number of keys in the right node
LFBLOC	CHARACTER*256	left node
RTBLOC	CHARACTER*256	right node
MKL	INTEGER*4	maximum key length
KEYVAL	CHARACTER*20	value of the last key examined
RTNOD	INTEGER*4	number of the right node
IO	BYTE	input/output device number
NAME	CHARACTER*9	name of the current tree
HTNOD	INTEGER*4	height of the node stack
MAXNOD	INTEGER*4	largest number yet assigned to a node
ONE	INTEGER*4	parameter equal to MKL + 2

OUTPUTS:

none

EXTERNALS:

STACK

SUBROUTINE ADJACENT

PURPOSE: To find a node adjacent to the current node.

INPUTS:

BUF	CHARACTER*256	array which contains the nodes in the current path
LEVEL	BYTE	level of the current node in the tree
NOD	INTEGER*4	number of the current node
BLOC	CHARACTER*256	current node
MANY	INTEGER*4	number of keys in the current node
ONE	INTEGER*4	parameter equal to MKL + 2

OUTPUTS:

RTNOD	INTEGER*4	number of the node to the right of the current node, if it exists; otherwise, the number of the current node
RTBLOC	CHARACTER*256	right node
RTM	INTEGER*4	number of keys in the right node
LFNOD	INTEGER*4	number of the current node, if no right node exists; otherwise, the number of the node to the left of the current node
LFBLOC	CHARACTER*245	left node
LFM	INTEGER*4	number of keys in the left node
KEYVAL	CHARACTER*20	value of the separator key

EXTERNALS:

VAL

SUBROUTINE DELKEY

PURPOSE: To delete a key and update the tree.

INPUTS:

IO	BYTE	input/output device number
NAME	CHARACTER*9	name of the current tree
HTREC	INTEGER*4	height of the record stack
MAXREC	INTEGER*4	largest number yet assigned to a data record
FULL	INTEGER*4	maximum number of keys in a node
HEIGHT	BYTE	height of the current tree
ONE	INTEGER*4	parameter equal to MKL + 2

OUTPUTS:

ROOT	INTEGER*4	node number of the root of the updated tree
HEIGHT	BYTE	height of the updated tree
HTREC	INTEGER*4	height of the updated record stack
MAXREC	INTEGER*4	largest number yet assigned to a data record
HTNOD	INTEGER*4	height of the updated node stack
MAXNOD	INTEGER*4	largest number yet assigned to a node

EXTERNALS:

GETKEY,STACK,ADJACENT,MERGE,PARENT,SHARE

SUBROUTINE SUCCESSOR

PURPOSE: To search for the key following the last key accessed.

INPUTS:

MARK	INTEGER*4	position in the current node of the last key examined
MANY	INTEGER*4	number of keys in the current node
ONE	INTEGER*4	parameter equal to MKL + 2
BLOC	CHARACTER*256	current node

OUTPUTS:

MARK	INTEGER*4	position of the next key in the node containing it
MANY	INTEGER*4	number of keys in the node containing the next key
NOD	INTEGER*4	number of the node containing the next key
PTR	INTEGER*4	number of the data record associated with the next key
KEYVAL	CHARACTER*20	value of the next key
ERR	BYTE	error code number

EXTERNALS:

none

SUBROUTINE FIRST

PURPOSE: To search for the first key in the tree.

INPUTS:

IO	BYTE	input/output device number
ROOT	INTEGER*4	node number of the tree root
HEIGHT	BYTE	height of the current tree

OUTPUTS:

MARK	INTEGER*4	position in the current node of the last key examined
MANY	INTEGER*4	number of keys in the current node
NOD	INTEGER*4	number of the current node
PTR	INTEGER*4	number of the data record associated with the last key examined
KEYVAL	CHARACTER*20	value of the last key examined
ERR	BYTE	error code number

EXTERNALS:

VAL

SUBROUTINE NEWROOT

PURPOSE: To create a new root.

INPUTS:

LFNOD	INTEGER*4	number of left node
RTNOD	INTEGER*4	number of right node
KEYVAL	CHARACTER*20	value of the first key in the right node
IO	BYTE	input/output device number
NAME	CHARACTER*9	name of the current tree
HTNOD	INTEGER*4	height of the node stack
MAXNOD	INTEGER*4	largest number yet assigned to a node

OUTPUTS:

ROOT	INTEGER*4	node number of the tree root
HEIGHT	BYTE	height of the current tree

EXTERNALS:

SYM,STACK

SUBROUTINE PARENT

PURPOSE: To update a parent node.

INPUTS:

LEVEL	BYTE	level of node last examined
LFNOD	INTEGER*4	number of left node
RTNOD	INTEGER*4	number of right node
BUF	CHARACTER*256	array which contains the nodes constituting the current path
ONE	INTEGER*4	parameter equal to MKL + 2
PATH	INTEGER*4	array which contains the node numbers defining the path from the root to the current node
INC	INTEGER*4	variable which determines the appropriate update action to be taken on the parent node
IO	BYTE	input/output device number
MKL	INTEGER*4	maximum key length
KEYVAL	CHARACTER*20	value of the separator key
PTR	INTEGER*4	number of the node to which the separator key points

OUTPUTS:

NOD	INTEGER*4	number of the parent node
MANY	INTEGER*4	updated number of keys in the parent node
BLOC	CHARACTER*256	updated parent node

EXTERNALS:

VAL,SYM

SUBROUTINE SPLIT

PURPOSE: To split a full node into two half-full nodes.

INPUTS:

IO	BYTE	input/output device number
NAME	CHARACTER*9	name of the current tree
HTNOD	INTEGER*4	height of the node stack
MAXNOD	INTEGER*4	largest number yet assigned to a node
FULL	INTEGER*4	maximum number of keys in a node
ONE	INTEGER*4	parameter equal to MKL + 2
BLOC	CHARACTER*256	current node
LEAF	LOGICAL	.TRUE. if and only if the current node is a leaf
NOD	INTEGER*4	number of the current node

OUTPUTS:

KEYVAL	CHARACTER*20	value of key in middle of current node
LFNOD	INTEGER*4	number of left node
RTNOD	INTEGER*4	number of right node

EXTERNALS:

SYM

SUBROUTINE ADDKEY

PURPOSE: To insert a key into the tree.

INPUTS:

ONE	INTEGER*4	parameter equal to MKL + 2
FULL	INTEGER*4	maximum number of keys in a node
IO	BYTE	input/output device number
NAME	CHARACTER*9	name of the current tree
HTREC	INTEGER*4	height of the record stack
MAXREC	INTEGER*4	largest number yet assigned to a data record

OUTPUTS:

ERR	BYTE	error code number
-----	------	-------------------

EXTERNALS:

GETKEY,STACK,SYM,PARENT,NEWROOT

SUBROUTINE LOOK

PURPOSE: To search a given node for the pointer to the next node in the path (if $H < \text{HEIGHT}$), or to search a leaf for the desired key (if $H = \text{HEIGHT}$).

INPUTS:

H	INTEGER*4	tree level of the current node
MANY	INTEGER*4	number of keys in the current node
ONE	INTEGER*4	parameter equal to MKL + 2
KEYVAL	CHARACTER*20	search string
BLOC	CHARACTER*256	current node
IO	BYTE	input/output device number
HEIGHT	BYTE	height of the current tree

OUTPUTS:

MARK	INTEGER*4	position in the current node of the last key examined
PTR	INTEGER*4	number of next node in path if $H < \text{HEIGHT}$; number of the data record associated with the last key examined if $H = \text{HEIGHT}$

EXTERNALS:

none

NSWC TR 85-54

7
EXTERNALS:

LOOK

SUBROUTINE GETKEY

PURPOSE: To search for the first occurrence of the key having its first MATCH characters equal to the value of KEYVAL.

INPUTS:

ABC	CHARACTER*1	command letter
KEYVAL	CHARACTER*20	search string
IO	BYTE	input/output device number
MKL	INTEGER*4	maximum key length
ROOT	INTEGER*4	node number of the tree root
HEIGHT	BYTE	height of the current tree

OUTPUTS:

MATCH	INTEGER*4	number of characters in the search string
PTR	INTEGER*4	number of the data record associated with the last key examined
NOD	INTEGER*4	number of the current node
PATH	INTEGER*4	array which contains the node numbers defining the path from the root to the current node
BLOC	CHARACTER*256	current node
BUF	CHARACTER*256	array which contains the nodes in the current path
MANY	INTEGER*4	number of keys in the current node
MARK	INTEGER*4	position in the current node of the last key examined
KEYVAL	CHARACTER*20	value of last key examined
ERR	BYTE	error code number

SUBROUTINE OLDTREE

PURPOSE: To open an existing tree called NAME and initialize the associated parameters.

INPUTS:

IO	BYTE	input/output device number
NAME	CHARACTER*9	name of the current tree

OUTPUTS:

HTREC	INTEGER*4	height of the record stack
MAXREC	INTEGER*4	largest number yet assigned to a data record
HTNOD	INTEGER*4	height of the node stack
MAXNOD	INTEGER*4	largest number yet assigned to a node
ROOT	INTEGER*4	node number of the tree root
HEIGHT	BYTE	height of the current tree

EXTERNALS:

none

SUBROUTINE NEWTREE

PURPOSE: To create a tree called NAME and initialize the associated parameters.

INPUTS:

IO	BYTE	input/output device number
NAME	CHARACTER*9	name of the current tree

OUTPUTS:

HTREC	INTEGER*4	height of the record stack
MAXREC	INTEGER*4	largest number yet assigned to a data record
HTNOD	INTEGER*4	height of the node stack
MAXNOD	INTEGER*4	largest number yet assigned to a node
ROOT	INTEGER*4	node number of the tree root
HEIGHT	BYTE	height of the current tree

EXTERNALS:

none

SUBROUTINE SHARE

PURPOSE: To equally redistribute keys between adjacent right and left nodes.

INPUTS:

LEAF	LOGICAL	.TRUE. if and only if the current node is a leaf
LFM	INTEGER*4	number of keys in the left node
RTM	INTEGER*4	number of keys in the right node
LFBLOC	CHARACTER*256	left node
KEYVAL	CHARACTER*20	value of separator key
RTBLOC	CHARACTER*256	right node
RTNOD	INTEGER*4	number of the right node
LFNOD	INTEGER*4	number of the left node
ONE	INTEGER*4	parameter equal to MKL + 2

OUTPUTS:

KEYVAL	CHARACTER*20	value of updated separator key
--------	--------------	--------------------------------

EXTERNALS:

SYM

SUBROUTINE STACK

PURPOSE: To push or pop either the node stack or the record stack, as required.

INPUTS:

PTR	INTEGER*4	number pushed onto the stack, if ADD = 1
IO	BYTE	input/output device number
NAME	CHARACTER*9	name of the current tree
EXT	CHARACTER*3	file name extender which determines which stack, node or record, is to be updated
ADD	BYTE	push/pop stack indicator
HT	INTEGER*4	height of the stack
MOST	INTEGER*4	largest number yet assigned to a node or data record, depending on which stack is to be updated

OUTPUTS:

HT	INTEGER*4	height of the updated stack
MOST	INTEGER*4	largest number yet assigned to a node (if EXT = 'NOD') or data record (if EXT = 'REC')
PTR	INTEGER*4	number popped from stack, if ADD = -1

EXTERNALS:

SYM,VAL

FUNCTION VAL

PURPOSE: To convert a 2-BYTE symbol into an integer between 0 and 65,635.
(See FUNCTION SYM for the inverse function.)

INPUTS:

A CHARACTER*2 2-BYTE symbol

OUTPUTS:

VAL INTEGER*4 integer between 0 and 65,635
corresponding to A

EXTERNALS:

none

FUNCTION SYM

PURPOSE: To convert an integer between 0 and 65,635
into a 2-BYTE symbol.
(See FUNCTION VAL for the inverse function.)

INPUTS:

NUM INTEGER*4 integer between 0 and 65,635

OUTPUTS:

SYM CHARACTER*2 2-BYTE symbol corresponding to NUM

EXTERNALS:

none

NSWC TR 85-54

PROGRAM DRIVER

PURPOSE:

To directly examine a tree by means of an interactive diagnostic program; especially effective when used in conjunction with a debug utility program.

INPUTS:

none

OUTPUTS:

none

EXTERNALS:

none

NSWC TR 85-54

APPENDIX B
FORTRAN CODE LISTING

SUBROUTINE BTREE(LTR,LDU,A,MAXLEN,IREC,IERR)

A 'B+ tree' is a data structure which is particularly well suited for storing the keys which identify the data records in a database. Data is rapidly retrieved by minimizing the number of mass storage accesses. This implementation allows a maximum of 65,535 keys, each having a maximum length of 20 characters; the tree has a maximum height of 5 and contains 256-byte nodes. Numbers corresponding to deleted nodes or records are placed in stacks and reused as needed. Application programs which call BTREE can have as many as 10 trees open simultaneously.

For a basic introduction to the subject of B+ trees, see "The Ubiquitous B Tree" by Douglas Comer, Computing Surveys, 11(1979)121-137; a more complete discussion can be found in "The Art of Computer Programming, Vol.3: Sorting and Searching" by Knuth, Addison-Wesley, 1973.

Complete documentation for BTREE is contained in "BTREE : A FORTRAN Code for a B+ Tree" by Elliot Winston, NSWC TR 85-54; a code for a database manager based on BTREE is the subject of "BOSS : A FORTRAN Code for a Relational Database Manager" by Elliot Winston, NSWC TR 85-56.

INPUTS:

LTR (CHARACTER*1)	COMMAND ACTION
A	(A)DD A KEY TO THE TREE
C	(C)REATE A NEW TREE
D	(D)ELETE A KEY FROM THE TREE
F	GET THE (F)IRST KEY IN THE TREE
G	(G)ET FIRST OCCURRENCE OF A TRUNCATED KEY
O	(O)PEN AN OLD TREE
S	GET THE (S)UCCESSOR KEY
LDU (BYTE)	UNIT NUMBER UNDER WHICH THE TREE COMMUNICATES WITH MASS STORAGE
A (CHARACTER*20)	KEY VALUE (LTR = 'A','D','G'); TREE NAME (LTR = 'C','O')
MAXLEN (INTEGER*4)	MAXIMUM KEY LENGTH (LTR = 'C')

OUTPUTS:

IREC (INTEGER*4)	NUMBER OF THE DATA RECORD ASSOCIATED WITH THE LAST KEY ACCESSED
A (CHARACTER*20)	VALUE OF LAST KEY ACCESSED

IERR (BYTE)	ERROR CODE
1	ILLEGAL VALUE OF PARAMETER 'LTR'
2	ATTEMPT TO CREATE EXISTENT TREE
3	ATTEMPT TO ACCESS NONEXISTENT TREE
4	ATTEMPT TO FIND NONEXISTENT KEY
5	NO SUCCESSOR EXISTS (LAST KEY IN TREE)
6	ATTEMPT TO INSERT KEY CURRENTLY IN TREE
IMPLICIT INTEGER*4 (A-Z)	
COMMON /XXXTREE/	
1	IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2	LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3	BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4	MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
	BYTE IO,ERR,LEVEL,HEIGHT
	CHARACTER ABC*1,NAME*9,KEYVAL*20
	CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
	BYTE IERR,LDU
	INTEGER MARKS(10),MANYS(10)
	CHARACTER LTR*1,A*20,BLOCS(10)*256
----- CONTROL LOGICAL FLOW OF A COMMAND -----	
	IO = LDU
	ABC = LTR
	ERR = 0
	IF (ABC.NE.'C'.OR.ABC.NE.'O') THEN
	ONE = MKL(IO) + 2
	FULL = 252/ONE
	END IF
	IF (ABC.EQ.'A') THEN
	KEYVAL = A
	CALL ADDKEY
	ELSE IF (ABC.EQ.'C') THEN
	NAME(IO) = A
	MKL(IO) = MAXLEN
	CALL NEWTREE
	ELSE IF (ABC.EQ.'D') THEN
	KEYVAL = A
	CALL DELKEY
	ELSE IF (ABC.EQ.'F') THEN
	CALL FIRST
	ELSE IF (ABC.EQ.'G') THEN
	KEYVAL = A
	CALL GETKEY
	ELSE IF (ABC.EQ.'0') THEN
	NAME(IO) = A

```
CALL OLDTREE
ELSE IF (ABC.EQ.'S') THEN
  MARK = MARKS(I0)
  MANY = MANYS(I0)
  BLOC = BLOCS(I0)
  CALL SUCCESSOR
ELSE
  ERR = 1
END IF
A = KEYVAL
IREC = PTR
IERR = ERR
IF (ABC.EQ.'G'.OR.ABC.EQ.'F'.OR.ABC.EQ.'S') THEN
  MARKS(I0) = MARK
  MANYS(I0) = MANY
  BLOCS(I0) = BLOC
END IF
RETURN
END
```

```

SUBROUTINE NEWTREE
C
IMPLICIT INTEGER*4 (A-Z)
C
COMMON /XXXTREE/
1   IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2   LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3   BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4   MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
     BYTE IO,ERR,LEVEL,HEIGHT
     CHARACTER ABC*1,NAME*9,KEYVAL*20
     CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
CHARACTER SYM*2,FN*13
LOGICAL*1 THERE
C
201 FORMAT(A256)
202 FORMAT(7I5)
C
C-----CREATE A NEW TREE
C-----CREATE A NEW TREE
C
CLOSE(UNIT=IO)
FN = NAME(IO)//'.KEY'
INQUIRE(FILE=FN,EXIST=THERE)
IF (THERE) THEN
  ERR = 2
  RETURN
END IF
OPEN(UNIT=IO,FILE=FN,STATUS='NEW',FORM='FORMATTED',
*           ACCESS='DIRECT',RECL=256)
I = 0
J = 0
BLOC(1:2) = SYM(I)
BLOC(3:4) = SYM(J)
WRITE(IO,201,REC=2) BLOC
J = 2
BLOC(3:4) = SYM(J)
WRITE(IO,201,REC=4) BLOC
HTREC(IO) = 0
MAXREC(IO) = 0
HTNOD(IO) = 0
MAXNOD(IO) = 2
ROOT(IO) = 2
HEIGHT(IO) = 0
C
ENTRY POINT FOR 'HEADER'
C
ENTRY HEADER
WRITE(IO,202,REC=1) HTREC(IO),MAXREC(IO),HTNOD(IO),MAXNOD(IO),
*           ROOT(IO),MKL(IO),HEIGHT(IO)

```

NSWC TR 85-54

RETURN
END

```

SUBROUTINE OLDTREE
C
IMPLICIT INTEGER*4 (A-Z)
C
COMMON /XXXTREE/
1   IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2   LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3   BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4   MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
BYTE IO,ERR,LEVEL,HEIGHT
CHARACTER ABC*1,NAME*9,KEYVAL*20
CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
CHARACTER*13 FN
LOGICAL*1 THERE
C
301 FORMAT(7I5)
C
C-----OPEN AND INITIALIZE AN OLD TREE
C-----
C
FN = NAME(IO)//'.KEY'
INQUIRE(FILE=FN,EXIST=THERE)
IF (THERE) THEN
  CLOSE(UNIT=IO)
  OPEN(UNIT=IO,FILE=FN,STATUS='OLD',FORM='FORMATTED',
*      ACCESS='DIRECT')
  READ(IO,301,REC=1) HTREC(IO),MAXREC(IO),HTNOD(IO),MAXNOD(IO),
*                  ROOT(IO),MKL(IO),HEIGHT(IO)
ELSE
  ERR = 3
  RETURN
END IF
RETURN
END

```

```

SUBROUTINE GETKEY
C      IMPLICIT INTEGER*4 (A-Z)
C
COMMON /XXXTREE/
1      IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2      LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3      BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4      MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
      BYTE IO,ERR,LEVEL,HEIGHT
      CHARACTER ABC*1,NAME*9,KEYVAL*20
      CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
401 FORMAT(A256)
C
C-----SEARCH FOR FIRST OCCURRENCE OF A KEY HAVING
C      FIRST 'MATCH' CHARACTERS EQUAL TO 'KEYVAL'
C-----
C
IF (ABC.EQ.'G') THEN
  J = 20
  DO WHILE (KEYVAL(J:J).EQ.' ')
    J = J - 1
  END DO
  MATCH = J
ELSE
  MATCH = MKL(IO)
END IF
PTR = ROOT(IO)
HEND = HEIGHT(IO)
DO 4015 H=0,HEND
  NOD = PTR
  PATH(H) = NOD
  READ(IO,401,REC=NOD) BLOC
  BUF(H) = BLOC
  MANY = VAL(BLOC(1:2))
  IF (MANY.EQ.0) THEN
C
    EMPTY TREE
C
    MARK = 1
    ERR = 4
    RETURN
  ELSE
    CALL LOOK(H)
  END IF
4015 CONTINUE
4020 IF (MATCH.EQ.MKL(IO).OR.MARK.LT.MANY+1) GO TO 4030
C
C      EXTENDED SEARCH FOR A STRICTLY TRUNCATED KEY
C

```

NSWC TR 85-54

```
J = MANY*ONE + 3
NOD = VAL(BLOC(J:J+1))
IF (NOD.EQ.0) GO TO 4030
PATH(HEND) = NOD
READ(10,401,REC=NOD) BLOC
BUF(HEND) = BLOC
MANY = VAL(BLOC(1:2))
CALL LOOK(HEND)
GO TO 4020
C
4030 K = (MARK-1)*ONE + 3
PTR = VAL(BLOC(K:K+1))
IF (KEYVAL.EQ.BLOC(K+2:K+MATCH+1)) THEN
  ERR = 0
  KEYVAL = BLOC(K+2:K+MKL(10)+1)
ELSE
  ERR = 4
END IF
RETURN
END
```



```

C      SUBROUTINE MERGE(LEAF)
C      IMPLICIT INTEGER*4 (A-Z)
C
C      COMMON /XXXTREE/
1      IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2      LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3      BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4      MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
      BYTE IO,ERR,LEVEL,HEIGHT
      CHARACTER ABC*1,NAME*9,KEYVAL*20
      CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
C      BYTE ADD
C      CHARACTER SYM*2,EXT*3
C      LOGICAL*1 LEAF
C
C      201 FORMAT(A256)
C
C-----C-----MERGE ADJACENT NODES INTO THE LEFT NODE-----C-----C
C
C      IF (LEAF) THEN
C          MANY = LFM + RTM
C          BLOC(1:2) = SYM(MANY)
C          BLOC(3:2+LFM*ONE) = LFBLOC(3:2+LFM*ONE)
C          I = 3 + LFM*ONE
C      ELSE
C          MANY = LFM + RTM + 1
C          BLOC(1:2) = SYM(MANY)
C          BLOC(3:4+LFM*ONE) = LFBLOC(3:4+LFM*ONE)
C          I = 5 + LFM*ONE
C          BLOC(I:I+MKL(IO)-1) = KEYVAL
C          I = I + MKL(IO)
C      END IF
C      BLOC(I:I+1+RTM*ONE) = RTBLOC(3:4+RTM*ONE)
C      WRITE(IO,201,REC=LFNOD) BLOC
C      ADD = 1
C      EXT = 'NOD'
C      CALL STACK(RTNOD,IO,NAME(IO),EXT,ADD,HTNOD(IO),MAXNOD(IO))
C      RETURN
CEND

```

```

SUBROUTINE ADJACENT
C
IMPLICIT INTEGER*4 (A-Z)
C
COMMON /XXXTREE/
1   IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2   LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3   BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4   MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
BYTE IO,ERR,LEVEL,HEIGHT
CHARACTER ABC*1,NAME*9,KEYVAL*20
CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
501 FORMAT(A256)
C
C-----FIND ADJACENT NODES
C-----C
C
MUCH = VAL(BUF(LEVEL-1)(1:2))
DO 5005 M=1,MUCH
   I = 3 + (M-1)*ONE
   IF (VAL(BUF(LEVEL-1)(I:I+1)).EQ.NOD) GO TO 5010
5005 CONTINUE
MARK = MUCH
I = 3 + (MARK-1)*ONE
RTNOD = NOD
RTBLOC = BLOC
RTM = MANY
LFNOD = VAL(BUF(LEVEL-1)(I:I+1))
READ(IO,501,REC=LFNOD) LFBLOC
LFM = VAL(LFBLOC(1:2))
GO TO 5015
5010 MARK = M
I = 3 + (MARK-1)*ONE
LFNOD = NOD
LFBLOC = BLOC
LFM = MANY
RTNOD = VAL(BUF(LEVEL-1)(I+ONE:I+1+ONE))
READ(IO,501,REC=RTNOD) RTBLOC
RTM = VAL(RTBLOC(1:2))
5015 KEYVAL = BUF(LEVEL-1)(I+2:I+1+MKL(IO))
RETURN
END

```

NSWC TR 85-54

```
IF (MANY.EQ.0) THEN
    ROOT(IO) = PATH(1)
    HEIGHT(IO) = HEIGHT(IO) - 1
    WRITE(IO,802,REC=1) HTREC(IO),MAXREC(IO),HTNOD(IO),
        MAXNOD(IO),ROOT(IO),MKL(IO),HEIGHT(IO)
*
END IF
WRITE(IO,801,REC=NOD) BLOC
RETURN
ELSE
    M = 2
    LEAF = .FALSE.
    GO TO 8005
END IF
ELSE
    CALL SHARE(LEAF)
    IF (MANY.EQ.LFM) RETURN
    INC = 0
    CALL PARENT(INC)
    WRITE(IO,801,REC=NOD) BLOC
    RETURN
END IF
END
```

```

SUBROUTINE DELKEY
C
IMPLICIT INTEGER*4 (A-Z)
C
COMMON /XXXTREE/
1   IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2   LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3   BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4   MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
BYTE IO,ERR,LEVEL,HEIGHT
CHARACTER ABC*1,NAME*9,KEYVAL*20
CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
BYTE ADD
CHARACTER SYM*2,EXT*3
LOGICAL*1 LEAF
C
801 FORMAT(A256)
802 FORMAT(7I5)
C
C-----.
C      DELETE A KEY FROM THE TREE
C-----.
C
C      DELETE A KEY FROM A LEAF
C
CALL GETKEY
IF (ERR.NE.0) RETURN
ADD = 1
EXT = 'REC'
CALL STACK(PTR,IO,NAME(IO),EXT,ADD,HTREC(IO),MAXREC(IO))
L = 3 + (MARK-1)*ONE
R = 4 + MANY*ONE
BLOC(L:R) = BLOC(L+ONE:R+ONE)
MANY = MANY - 1
BLOC(1:2) = SYM(MANY)
IF (MANY.GE.FULL/2.OR.HEIGHT(IO).EQ.0) THEN
    WRITE(IO,801,REC=NOD) BLOC
    RETURN
END IF
C
C      UPDATE TREE
C
LEVEL = HEIGHT(IO)
M = 1
LEAF = .TRUE.
8005 CALL ADJACENT
IF (LFM+RTM.LE.FULL-M) THEN
    CALL MERGE(LEAF)
    INC = -1
    CALL PARENT(INC)
    IF (MANY.GE.FULL/2.OR.LEVEL.EQ.0) THEN

```

```

SUBROUTINE SUCCESSOR
C
IMPLICIT INTEGER*4 (A-Z)
C
COMMON /XXXTREE/
1  IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2  LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3  BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4  MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
BYTE IO,ERR,LEVEL,HEIGHT
CHARACTER ABC*1,NAME*9,KEYVAL*20
CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
801 FORMAT(A256)
C
C-----.
C      GET THE NEXT KEY IN SEQUENCE FOLLOWING THE LAST
C      KEY ACCESSED UNDER THE SAME INPUT/OUTPUT NUMBER
C-----.
C
IF (MARK.LT.MANY) THEN
  MARK = MARK + 1
  GO TO 8010
ELSE
  I = 3 + MANY*ONE
  NOD = VAL(BLOC(I:I+1))
  IF (NOD.EQ.0) THEN
    ERR = 5
    RETURN
  ELSE
    READ(IO,801,REC=NOD) BLOC
    MANY = VAL(BLOC(1:2))
    MARK = 1
  END IF
END IF
8010 K = 3 + (MARK-1)*ONE
PTR = VAL(BLOC(K:K+1))
KEYVAL = BLOC(K+2:K+MKL(IO)+1)
RETURN
END

```

```

SUBROUTINE FIRST
C
IMPLICIT INTEGER*4 (A-Z)
C
COMMON /XXXTREE/
1   IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2   LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3   BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4   MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
BYTE IO,ERR,LEVEL,HEIGHT
CHARACTER ABC*1,NAME*9,KEYVAL*20
CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
701 FORMAT(A256)
C
C-----SEARCH FOR THE FIRST KEY IN THE TREE
C-----
C
NOD = ROOT(IO)
HEND = HEIGHT(IO)
DO 7005 H=0,HEND
  READ(IO,701,REC=NOD) BLOC
  IF (H.EQ.HEND) GO TO 7010
  NOD = VAL(BLOC(3:4))
7005 CONTINUE
7010 MARK = 1
  MANY = VAL(BLOC(1:2))
C
EMPTY TREE
C
IF (MANY.EQ.0) THEN
  ERR = 4
  RETURN
END IF
PTR = VAL(BLOC(3:4))
KEYVAL = BLOC(5:4+MKL(IO))
RETURN
END

```

```

C      SUBROUTINE NEWROOT
C      IMPLICIT INTEGER*4 (A-Z)
C
C      COMMON /XXXTREE/
1      IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2      LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3      BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4      MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
      BYTE IO,ERR,LEVEL,HEIGHT
      CHARACTER ABC*1,NAME*9,KEYVAL*20
      CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
C      BYTE ADD
C      CHARACTER SYM*2,EXT*3
C
C      701 FORMAT(A256)
C      702 FORMAT(7I5)
C
C-----CREATE A NEW ROOT
C-----
C
C      MANY = 1
BLOC(1:2) = SYM(MANY)
BLOC(3:4) = SYM(LFNOD)
BLOC(5:2+ONE) = KEYVAL
BLOC(3+ONE:4+ONE) = SYM(RTNOD)
ADD = -1
EXT = 'NOD'
CALL STACK(PTR,IO,NAME(IO),EXT,ADD,HTNOD(IO),MAXNOD(IO))
WRITE(IO,701,REC=PTR) BLOC
ROOT(IO) = PTR
HEIGHT(IO) = HEIGHT(IO) + 1
WRITE(IO,702,REC=1) HTREC(IO),MAXREC(IO),HTNOD(IO),MAXNOD(IO),
*                           ROOT(IO),MKL(IO),HEIGHT(IO)
RETURN
END

```

NSWC TR 85-54

BLOC(1:2) = SYM(MANY)
RETURN
END

```

      SUBROUTINE PARENT(INC)
C
      IMPLICIT INTEGER*4 (A-Z)
C
      COMMON /XXXTREE/
1      IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2      LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3      BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4      MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
      BYTE IO,ERR,LEVEL,HEIGHT
      CHARACTER ABC*1,NAME*9,KEYVAL*20
      CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
      CHARACTER*2 SYM
C
C----- UPDATE A PARENT NODE
C----- -----
C
      LEVEL = LEVEL - 1
      NOD = PATH(LEVEL)
      BLOC = BUF(LEVEL)
      MANY = VAL(BLOC(1:2))
      DO 6005 M=1,MANY
         I = 3 + (M-1)*ONE
         IF(VAL(BLOC(I:I+1)).EQ.LFNOD) GO TO 6010
6005 CONTINUE
      M = MANY + 1
6010 L = 5 + (M-1)*ONE
      R = 4 + MANY*ONE
      IF (INC.EQ.-1) THEN
C
      DELETE SEPARATOR FROM PARENT NODE
C
      IF (M.LT.MANY) BLOC(L:R) = BLOC(L+ONE:R+ONE)
C
      ELSE IF (INC.EQ.0) THEN
C
      UPDATE VALUE OF SEPARATOR IN PARENT NODE
C
      BLOC(L:L+MKL(IO)-1) = KEYVAL
C
      ELSE
C
      INSERT SEPARATOR INTO PARENT NODE
C
      IF (M.LE.MANY) BLOC(L+ONE:R+ONE) = BLOC(L:R)
      BLOC(L:L+MKL(IO)-1) = KEYVAL
      BLOC(L+MKL(IO):L+ONE-1) = SYM(RTNOD)
C
      END IF
      MANY = MANY + INC

```

```

SUBROUTINE SPLIT(LEAF)
C
IMPLICIT INTEGER*4 (A-Z)
C
COMMON /XXXTREE/
1   IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2   LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3   BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4   MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
BYTE IO,ERR,LEVEL,HEIGHT
CHARACTER ABC*1,NAME*9,KEYVAL*20
CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
BYTE ADD
CHARACTER SYM*2,EXT*3
LOGICAL*1 LEAF
C
601 FORMAT(A256)
C
C-----SPLIT A FULL NODE INTO TWO HALF-FULL NODES
C-----
C
ADD = -1
EXT = 'NOD'
CALL STACK(PTR,IO,NAME(IO),EXT,ADD,HTNOD(IO),MAXNOD(IO))
MANY = FULL/2
LFNOD = NOD
LFBLOC(1:2) = SYM(MANY)
I = 2 + MANY*ONE
KEYVAL = BLOC(I+3:I+ONE)
IF (LEAF) THEN
  LFBLOC(3:I) = BLOC(3:I)
  LFBLOC(I+1:I+2) = SYM(PTR)
  WRITE(IO,601,REC=NOD) LFBLOC
  MANY = FULL - MANY
ELSE
  LFBLOC(3:I+2) = BLOC(3:I+2)
  WRITE(IO,601,REC=NOD) LFBLOC
  I = I + ONE
  MANY = FULL - 1 - MANY
END IF
RTNOD = PTR
RTBLOC(1:2) = SYM(MANY)
RTBLOC(3:4+MANY*ONE) = BLOC(I+1:4+FULL*ONE)
WRITE(IO,601,REC=PTR) RTBLOC
RETURN
END

```

```
IF (LEVEL.GT.0) THEN
  INC = 1
  CALL PARENT(INC)
  IF (MANY.LT.FULL) THEN
    WRITE(10,401,REC=NOD) BLOC
    RETURN
  ELSE
    LEAF = .FALSE.
    GO TO 4005
  END IF
ELSE
  CALL NEWROOT
END IF
RETURN
END
```

```

SUBROUTINE ADDKEY
C
IMPLICIT INTEGER*4 (A-Z)
C
COMMON /XXXTREE/
1   IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2   LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3   BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4   MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
BYTE IO,ERR,LEVEL,HEIGHT
CHARACTER ABC*1,NAME*9,KEYVAL*20
CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
BYTE ADD
CHARACTER SYM*2,EXT*3
LOGICAL*1 LEAF
C
401 FORMAT(A256)
C
C-----  

C      INSERT A KEY INTO THE TREE  

C-----  

C
C      INSERT A KEY INTO A LEAF
C
CALL GETKEY
IF (ERR.EQ.0) THEN
  ERR = 6
  RETURN
ELSE
  ERR = 0
END IF
ADD = -1
EXT= 'REC'
CALL STACK(PTR,IO,NAME(IO),EXT,ADD,HTREC(IO),MAXREC(IO))
L = 3 + (MARK-1)*ONE
R = 4 + MANY*ONE
BLOC(L+ONE:R+ONE) = BLOC(L:R)
BLOC(L:L+1) = SYM(PTR)
BLOC(L+2:L+ONE-1) = KEYVAL
MANY = MANY + 1
BLOC(1:2) = SYM(MANY)
IF (MANY.LT.FULL) THEN
  WRITE(IO,401,REC=NOD)BLOC
  RETURN
END IF
C
C      UPDATE TREE
C
LEVEL = HEIGHT(IO)
LEAF = .TRUE.
4005 CALL SPLIT(LEAF)

```

```

C      SUBROUTINE LOOK(H)
C      IMPLICIT INTEGER*4 (A-Z)
C
C      COMMON /XXXTREE/
1      IO,ERR,MARK,MANY,PTR,NOD,LEVEL,MATCH,ONE,FULL,ABC,
2      LFM,RTM,LFNOD,RTNOD,KEYVAL,BLOC,LFBLOC,RTBLOC,
3      BUF(0:5),PATH(0:5),HTREC(10),MAXREC(10),HTNOD(10),
4      MAXNOD(10),ROOT(10),MKL(10),HEIGHT(10),NAME(10)
        BYTE IO,ERR,LEVEL,HEIGHT
        CHARACTER ABC*1,NAME*9,KEYVAL*20
        CHARACTER*256 BUF,BLOC,LFBLOC,RTBLOC
C
C      CHARACTER*20 TRY
C
C-----C      SEARCH A GIVEN NODE
C-----C
C      DO 5015 M=1,MANY
        K = 5 + (M-1)*ONE
        TRY = BLOC(K:K+MKL(IO)-1)
        IF (TRY.GE.KEYVAL) GO TO 5020
5015 CONTINUE
        MARK = MANY + 1
        GO TO 5025
5020 IF (H.EQ.HEIGHT(IO)) THEN
        MARK = M
    ELSE
        IF (TRY.EQ.KEYVAL) THEN
            MARK = M + 1
        ELSE
            MARK = M
        END IF
    END IF
5025 K = (MARK-1)*ONE + 3
        PTR = VAL(BLOC(K:K+1))
        RETURN
    END

```

```
J = 2 + (MANY-LFM)*ONE
BLOC(I+1:I+J-2) = RTBLOC(3:J)
K = I + J - 1
BLOC(K:K+1) = SYM(RTNOD)
ELSE
  BLOC(3:I+2) = LFBLOC(3:I+2)
  BLOC(I+3:I+ONE) = KEYVAL
  J = 2 + (MANY-1-LFM)*ONE
  I = I + ONE
  BLOC(I+1:I+J) = RTBLOC(3:J+2)
END IF
WRITE(10,201,REC=LFNOD) BLOC
K = 3 + (MANY - LFM)*ONE
MANY = LFM + RTM - MANY
BLOC(3:4+MANY*ONE) = RTBLOC(K:4+RTM*ONE)
KEYVAL = RTBLOC(J+3:J+ONE)
END IF
BLOC(1:2) = SYM(MANY)
WRITE(10,201,REC=RTNOD) BLOC
RETURN
END
```

```

SUBROUTINE STACK(PTR,IO,NAME,EXT,ADD,HT,MOST)
C
IMPLICIT INTEGER*4 (A-Z)
C
CHARACTER SYM*2,EXT*3,NAME*9,FN*13,B*128
BYTE ADD,IO
C
901 FORMAT(A128)
C
C-----*
C      PUSH/POP (ACCORDING TO ADD = 1,-1) THE RECORD/NODE
C      STACK (ACCORDING TO EXT = 'REC','NOD')
C-----*
C
IF (HT.GT.0.OR.ADD.EQ.1) THEN
  CLOSE(UNIT=IO)
  FN = NAME//'.'//EXT
  OPEN(UNIT=IO,FILE=FN,STATUS='UNKNOWN',FORM='FORMATTED',
*          ACCESS='DIRECT',RECL=128)
  I = 1 + HT/64
  J = MOD(HT,64)
  K = 2*j
  HT = HT + ADD
  IF (ADD.EQ.1) THEN
    IF (K.GT.0) READ(IO,901,REC=I) B
    B(K+1:K+2) = SYM(PTR)
    WRITE(IO,901,REC=I) B
  ELSE
    IF (J.EQ.0) THEN
      I = I - 1
      K = 126
    ELSE
      K = K - 2
    END IF
    READ(IO,901,REC=I) B
    PTR = VAL(B(K+1:K+2))
  END IF
  IF (HT.EQ.0) THEN
    CLOSE(UNIT=IO,STATUS='DELETE')
  ELSE
    CLOSE(UNIT=IO)
  END IF
  FN = NAME//'.KEY'
  OPEN(UNIT=IO,FILE=FN,STATUS='OLD',FORM='FORMATTED',
*          ACCESS='DIRECT')
  ELSE
    MOST = MOST + 1
    PTR = MOST
  END IF
C
ENTRY POINT OF 'HEADER' IS LOCATED IN 'NEWTREE'

```

NSWC TR 85-54

CALL HEADER
RETURN
END

```
C      FUNCTION VAL(A)
C      IMPLICIT INTEGER*4 (A-Z)
C      CHARACTER*2 A
C-----
C      CONVERT CODED 2-BYTE SYMBOL INTO AN
C      INTEGER BETWEEN 0 AND 65535
C-----
C
VAL = ICHAR(A(1:1))
J = ICHAR(A(2:2))
K = 2**8
DO 1005 I=0,7
    VAL = VAL + IBITS(J,I,1)*K
    K = 2*K
1005 CONTINUE
RETURN
END
```

```
CHARACTER*2 FUNCTION SYM(NUM)
C
C      IMPLICIT INTEGER*4 (A-Z)
C
C-----  
C      CONVERT AN INTEGER BETWEEN 0 AND 65535
C      INTO A CODED 2-BYTE SYMBOL
C-----  
C
DO 2010 J=1,2
      SUM = 0
      K = 1
      DO 2005 I=0,7
            SUM = SUM + IBITS(NUM,(J-1)*8+I,1)*K
            K = 2*K
2005    CONTINUE
      SYM(J:J) = CHAR(SUM)
2010    CONTINUE
      RETURN
      END
```

```

PROGRAM DRIVER
C
CHARACTER LTR*1,A*20
C
10 FORMAT(A1)
11 FORMAT(A25)
21 FORMAT(10X,' ILLEGAL VALUE OF LTR')
22 FORMAT(10X,' TREE CURRENTLY EXISTS')
23 FORMAT(10X,' NONEXISTENT TREE')
24 FORMAT(10X,' CANNOT FIND DESIRED KEY VALUE')
25 FORMAT(10X,' NO SUCCESSOR KEY EXISTS')
26 FORMAT(10X,' KEY CURRENTLY EXISTS IN TREE')

C
C-----  

C      PROGRAM TO DIRECTLY EXAMINE A TREE
C-----  

C
      WRITE(6,*)' ENTER LOGICAL UNIT NUMBER'
      READ(5,*) LDU
100  WRITE(6,*) ' '
      WRITE(6,*)' A - add      D - delete    G - get      key'
      WRITE(6,*)'          F - first     S - successor   key'
      WRITE(6,*)'          O - open      C - create     tree'
      WRITE(6,*)' ENTER LETTER'
      READ(5,10) LTR
      IF (LTR.EQ.'A'.OR.LTR.EQ.'D'.OR.LTR.EQ.'G') THEN
        WRITE(6,*)' ENTER KEY VALUE'
        READ(5,11) A
      ELSE IF (LTR.EQ.'O'.OR.LTR.EQ.'C') THEN
        WRITE(6,*)' ENTER TREE NAME'
        READ(5,11) A
        IF (LTR.EQ.'C') THEN
          WRITE(6,*)' ENTER LENGTH OF PRIMARY KEY'
          READ(5,*) MAXLEN
        END IF
      END IF
      CALL BTREE(LTR,LDU,A,MAXLEN,IREC,IERR)

      IF (IERR.EQ.0) THEN
        GO TO 200
      ELSE IF (IERR.EQ.1) THEN      ! ILLEGAL VALUE OF 'LTR'
        WRITE(6,21)
      ELSE IF (IERR.EQ.2) THEN      ! TREE CURRENTLY EXISTS
        WRITE(6,22)
      ELSE IF (IERR.EQ.3) THEN      ! NONEXISTENT TREE
        WRITE(6,23)
      ELSE IF (IERR.EQ.4) THEN      ! CANNOT FIND KEY
        WRITE(6,24)
      ELSE IF (IERR.EQ.5) THEN      ! NO SUCCESSOR KEY
        WRITE(6,25)
      ELSE IF (IERR.EQ.6) THEN      ! KEY CURRENTLY EXISTS
        WRITE(6,26)

```

```
END IF
WRITE(6,*)
WRITE(6,*)
      REQUEST VOIDED'
WRITE(6,*)

C 200 WRITE(6,*)
      DO YOU WISH TO EXIT? (Y/N)'
      READ (5,10) LTR
      IF(LTR.EQ.'N') GO TO 100
      STOP
      END
```

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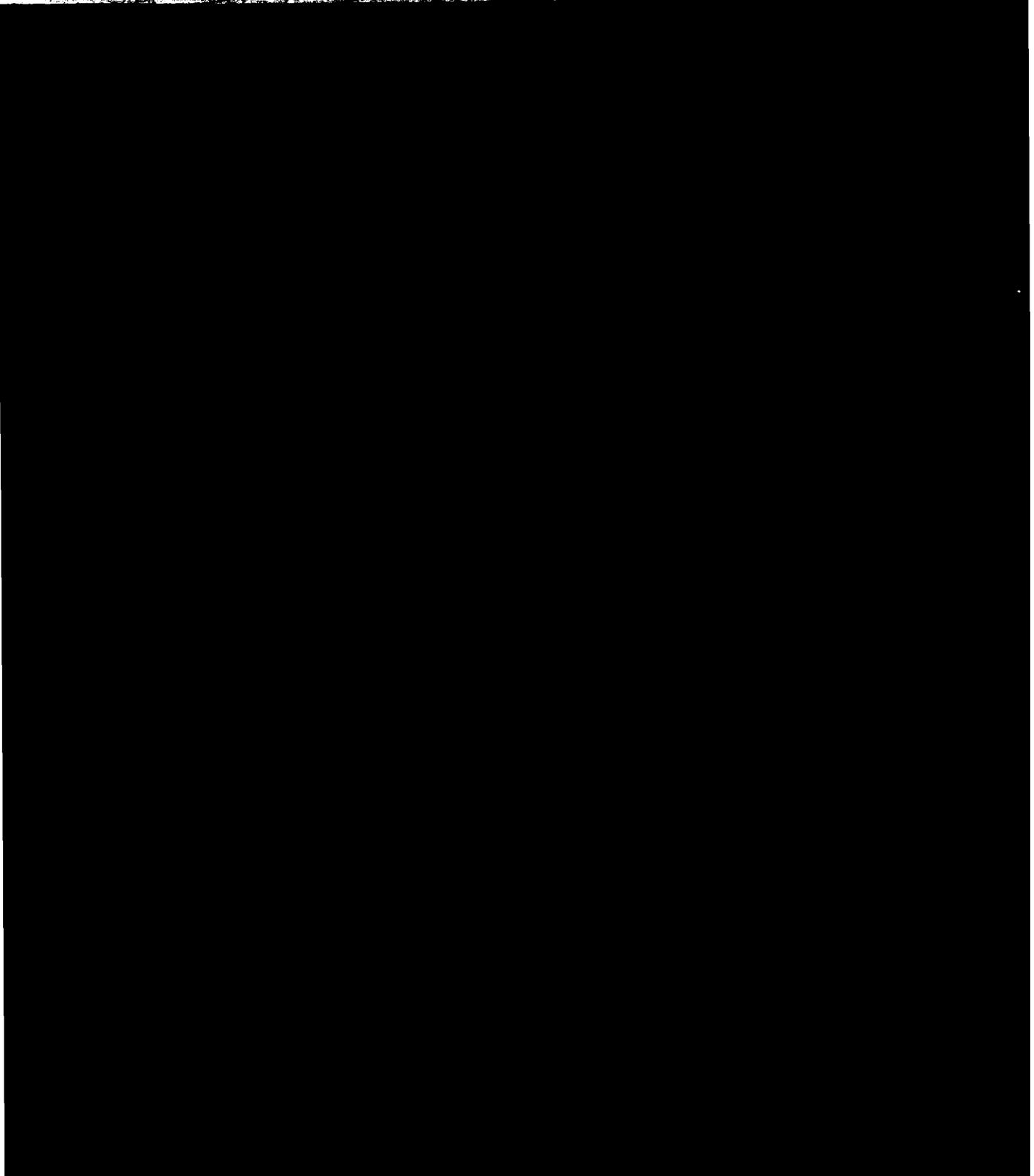
C
LEVEL = HEIGHT(I0)
M = 1
LEAF = .TRUE.
8005 CALL ADJACENT
IF (LFM+RTM.LE.FULL-M) THEN
CALL MERGE(LEAF)
INC = -1
CALL PARENT(INC)
IF (MANY.GE.FULL/2.OR.LEVEL.EQ.0) THEN

B-19

B-20

RETURN
END

B-21



END

B-22

```
BLOC(3:K) = LFBLOC(I+ONE+1:J)
BLOC(K+I:K+MKL(IO)) = KEYVAL
END IF
K = 3 + (LFM - MANY)*ONE
MANY = LFM + RTM - MANY
BLOC(K:4+MANY*ONE) = RTBLOC(3:4+RTM*ONE)
KEYVAL = LFBLOC(I+3:I+ONE)
ELSE
I = 2 + LFM*ONE
IF (LEAF) THEN
    BLOC(3:I) = LFBLOC(3:I)
```

B-24

```
*      ACCESS='DIRECT')
ELSE
  MOST = MOST + 1
  PTR = MOST
END IF
C
C
C      ENTRY POINT OF 'HEADER' IS LOCATED IN 'NEWTREE'
```

B-26

B-27

B-28

```
      WRITE(6,22)
ELSE IF (IERR.EQ.3) THEN          ! NONEXISTENT TREE
      WRITE(6,23)
ELSE IF (IERR.EQ.4) THEN          ! CANNOT FIND KEY
      WRITE(6,24)
ELSE IF (IERR.EQ.5) THEN          ! NO SUCCESSOR KEY
      WRITE(6,25)
ELSE IF (IERR.EQ.6) THEN          ! KEY CURRENTLY EXISTS
      WRITE(6,26)
```

B-30

(1/2)

